Enzo-P / Cello

Formation of the First Galaxies

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Blue Waters Symposium 2013
National Center for Supercomputing Applications
University of Illinois at Urbana-Champaign

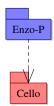
Introducing Enzo-P / Cello

Our group actively develops two related parallel applications:



Enzo: astrophysics / cosmology application

- patch-based adaptive mesh refinement (AMR)
- MPI or MPI/OpenMP
- almost 20 years development



Enzo-P / Cello: "Petascale" fork of Enzo code

- "forest of octrees" AMR
- Charm++ or MPI
- \bullet ≈ 3 years development
- work in progress–AMR just coming online

Enzo's strengths

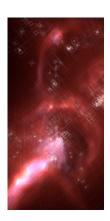


[John Wise]

- Spans multiple application domains
 - astrophysical fluid dynamics
 - hydrodynamic cosmology
- Rich multi-physics capabilities
 - fluid, particle, gravity, radiation, ...
- Extreme resolution range
 - 34 levels of refinement by 2!
- Active global development community
 - ≈ 25 developers

Enzo's struggles

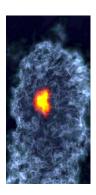
- Memory usage
 - ≈ 1.5KB/patch (MPI/OpenMP helps)
 - memory fragmentation
- Mesh quality
 - 2-to-1 constraint can be violated
 - asymmetric mesh for symmetric problem
- Load balancing
 - difficulty maintaining parent-child locality
- Parallel scaling
 - AMR overhead dominates computation



[Tom Abel, John Wise, Ralf Kaehler]

Enzo's pursuit of scalability

- Enzo was born in early 1990's
- "Extreme" meant 100 processors
- Continual scalability improvements
 - MPI/OpenMP parallelism
 - "neighbor-finding" algorithm
 - I/O optimizations
- Further improvement getting harder
 - increasing scalability requirements
 - easy improvements made already
- Motivates concurrent rewriting
 - Enzo-P "Petascale" Enzo fork
 - Cello AMR framework



[Sam Skillman, Matt Turk]

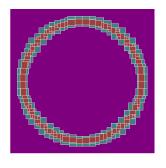
Enzo-P / Cello design overview



- Charm++ parallelism
 - asynchronous, data-driven
 - latency tolerant
 - dynamic load balancing
 - checkpoint / restart

- Octree-based AMR
 - "forest" for root mesh
 - easier to implement
 - scalability advantages
 - fast neighbor-finding

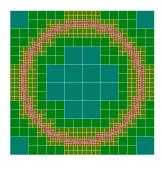
Some advantages of patch-based AMR



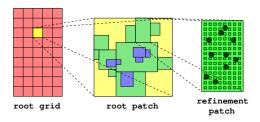
- Flexible patch size and shape
 - improved refinement efficiency
- Larger patches
 - smaller surface/volume ratio
 - reduced communication
 - amortized loop overhead
- Fewer patches
 - reduced AMR meta-data

Some advantages of octree-based AMR

- Fixed block size and shape
 - simplified load balancing
 - dynamic memory reuse
- More blocks
 - more parallelism available
- Smaller nodes
 - reduced AMR meta-data
- Compute only on leaf nodes
 - less communication

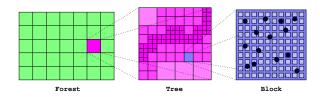


Enzo's AMR data structure



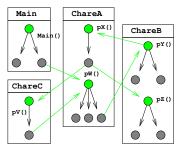
- Patches assigned to MPI processes
- Refinement patches created on root patch process
- Load balancing relocates refinement patches
- Patch data (grid, particle) are distributed
- Replicated AMR hierarchy structure

Enzo-P / Cello's AMR data structure



- Each block is a Charm++ chare
- Blocks initially mapped to root node process
- Charm++ load balances
- AMR hierarchy structure is fully distributed

Charm++ program structure



A Charm++ Program

- Charm++ program
 - Charm++ objects are *chares*
 - invoke remote *entry methods*
 - communicate via messages
- Charm++ runtime system
 - schedules entry methods
 - maps chares to processors
 - migrates chares to balance
- Additional scalability features
 - checkpoint / restart
 - sophisticated DLB strategies

Charm++ collections of chares

Chare Array



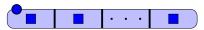
- distributed array of chares
- migrateable elements
- flexible indexing

Chare Group



• one chare per processor (non-migrateable)

Chare Nodegroup



one chare per node (non-migrateable)

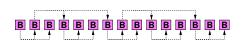
Cello implementation options using Charm++

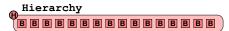
1. Singleton chares

- unlimited hierarchy depth
- tedious to program
- limited Charm++ support

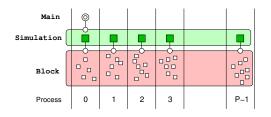
2. Chare array

- efficient: single access
- restricted hierarchy depth





Charm++ entities in Enzo-P / Cello

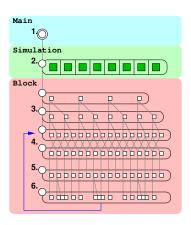


- "mainchare" called at program startup
- Simulation chare group holds global data
- Block chare array defines forest of octrees

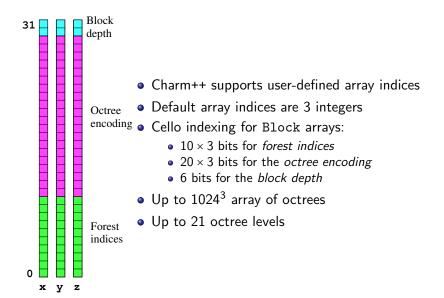
Control flow in Enzo-P / Cello

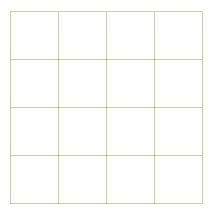
Current Enzo-P / Cello control flow

- Startup
- Initialize
- Mesh creation
- Ghost refresh
- Computation
- Mesh adaptation

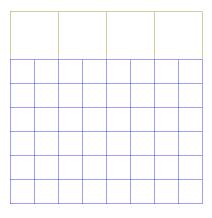


Block chare array indexing

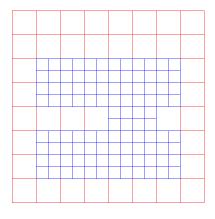




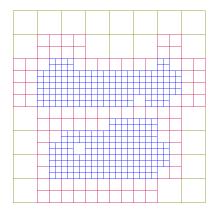
- Begins with the forest root grid
- Proceeds level-by-level
- Blocks evaluate refinement criteria
 - if refine, create child blocks
 - if coarsen, notify parent block
- Refine can violate 2-1 constraint
 - tell coarse neighbors to refine
 - may recurse
- Quiescence detection between steps
- Keep track of neighbors and children



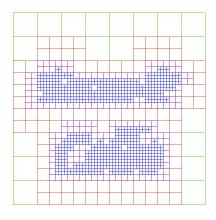
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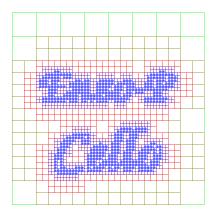
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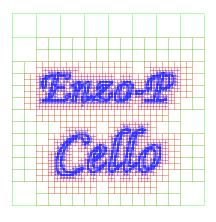
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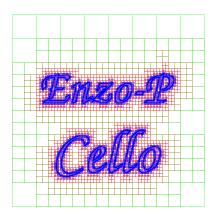
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Cello AMR ghost zone refresh

Intra-level refresh

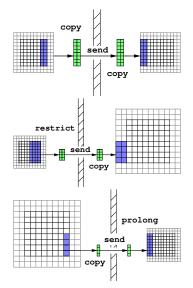
- 1. FaceBlock loads face cells
- 2. Charm++ entry method send
- 3. FaceBlock stores ghost cells

Fine-to-coarse refresh

- 1. FaceBlock coarsens face cells
- 2. Charm++ entry method send
- 3. FaceBlock stores ghost cells

Coarse-to-fine refresh

- 1. FaceBlock loads face cells
- 2. Charm++ entry method send
- 3. FaceBlock interpolates ghost cells



Summary

	Enzo	Enzo-P / Cello
Parallelization	MPI/OpenMP	Charm++
AMR	patch-based	tree-based
AMR structure	replicated	distributed
Block sizes	×1000 variation	constant
Task scheduling	level-parallel	dependency-driven
Load balancing	patch migration	Charm++
Fault tolerance	checkpoint/restart	Charm++

http://cello-project.org

NSF PHY-1104819, AST-0808184